

COL7160 : Quantum Computing

Lecture 19: Grover's Search Algorithm and Amplitude Amplification

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1 Definitions

Definition 1 (Search Problem). Given a function $f : \{0, 1\}^n \rightarrow \{0, 1\}$, find a x such that $f(x) = 1$

Definition 2. Lets denote the following notation for convinience

- Domain U : which in this case is $\{0, 1\}^n$
- Total number of elements : $N = |U| = 2^n$
- Set of tagged elements : $A = \{x \in U \mid f(x) = 1\}$
- Number of tagged elements : $t = |A|$
- Complement of A : $B = A^c = \{x \in U \mid f(x) = 0\}$

Classically, if you pick an element at random from U , the probability of it being tagged is $|A|/|U| = t/N$, so the expected number of elements we'd need to sample would be N/t

2 Quantum Setting

Definition 3 (Function Oracle). Access to the function f is provided as an oracle Z_f defined by its action on the basis states

$$Z_f |x\rangle = (-1)^{f(x)} |x\rangle \quad (1)$$

Definition 4 (Uniform Superposition State). The uniform superposition state on a set S is $|S\rangle$ such that

$$|S\rangle = \frac{1}{\sqrt{|S|}} \sum_{x \in S} |x\rangle \quad (2)$$

Lemma 5 (Uniform State). The uniform superposition state over the domain $U = \{0, 1\}^n$ is generated by applying n Hadamard gates in parallel to $|0^n\rangle$

$$|U\rangle = H^{\otimes n} |0^n\rangle \quad (3)$$

Proof. The action of a single Hadamard gate on the basis state $|0\rangle$ is given by $H|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$. Applying n such gates in parallel to $|0^n\rangle = |0\rangle \otimes \dots \otimes |0\rangle$ yields:

$$H^{\otimes n} |0^n\rangle = \bigotimes_{i=1}^n H|0\rangle = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}} \right)^{\otimes n} = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle = |U\rangle \quad (4)$$

□

Definition 6 (OR Function).

$$\text{OR}(x) = \begin{cases} 0 & \text{if } x = 0 \\ 1 & \text{otherwise} \end{cases} \quad (5)$$

Lemma 7 (OR Oracle). The oracle for OR, acting as defined in 3 can be written as

$$Z_{\text{OR}} = 2|0\rangle\langle 0| - I \quad (6)$$

Proof. By Definition 3, the oracle Z_{OR} can be expanded in the computational basis as:

$$Z_{\text{OR}} = \sum_{x \in \{0,1\}^n} (-1)^{\text{OR}(x)} |x\rangle \langle x| \quad (7)$$

Splitting the sum into the $x = 0$ and $x \neq 0$ terms, and using the definition of the OR function:

$$Z_{\text{OR}} = (-1)^0 |0\rangle \langle 0| + \sum_{x \neq 0} (-1)^1 |x\rangle \langle x| = |0\rangle \langle 0| - \sum_{x \neq 0} |x\rangle \langle x| \quad (8)$$

Using the completeness relation $I = |0\rangle \langle 0| + \sum_{x \neq 0} |x\rangle \langle x|$, we can substitute $\sum_{x \neq 0} |x\rangle \langle x| = I - |0\rangle \langle 0|$:

$$Z_{\text{OR}} = |0\rangle \langle 0| - (I - |0\rangle \langle 0|) = 2|0\rangle \langle 0| - I \quad (9)$$

□

Lemma 8 (Reflections). *Any operator of the form $Z = 2|\psi\rangle \langle \psi| - I$ is a reflection about $|\psi\rangle$*

Proof. Using the resolution of identity $I = |\psi\rangle \langle \psi| + \sum_i |\psi_i^\perp\rangle \langle \psi_i^\perp|$, where $\{|\psi_i^\perp\rangle\}$ is an orthonormal basis for the subspace orthogonal to $|\psi\rangle$:

$$Z = 2|\psi\rangle \langle \psi| - \left(|\psi\rangle \langle \psi| + \sum_i |\psi_i^\perp\rangle \langle \psi_i^\perp| \right) = |\psi\rangle \langle \psi| - \sum_i |\psi_i^\perp\rangle \langle \psi_i^\perp| \quad (10)$$

This operator preserves the component along $|\psi\rangle$ and negates all components in the orthogonal subspace, defining a reflection about $|\psi\rangle$. □

Corollary 9 (Reflection about U). *The operator $H^{\otimes n} (2|0\rangle \langle 0| - I) H^{\otimes n}$ is a reflection about the uniform state U .*

Proof. Let $Z = 2|0\rangle \langle 0| - I$ be the reflection about $|0\rangle$ (as in Lemma 7). Applying the similarity transformation $H^{\otimes n}$ to Z :

$$H^{\otimes n} Z H^{\otimes n} = H^{\otimes n} (2|0\rangle \langle 0| - I) H^{\otimes n} = 2(H^{\otimes n} |0\rangle) (\langle 0| H^{\otimes n}) - H^{\otimes n} I H^{\otimes n} \quad (11)$$

By Lemma 5, $H^{\otimes n} |0\rangle = |U\rangle$. Since $H^{\otimes n}$ is unitary and self-adjoint ($H = H^\dagger$, $H^2 = I$), we have:

$$2|U\rangle \langle U| - I \quad (12)$$

By Lemma 8, this operator is a reflection about $|U\rangle$. □

Proposition 10. *The uniform superposition can be written as a combination of the uniform superposition states over A, B as:*

$$|U\rangle = \sqrt{\frac{t}{N}} |A\rangle + \sqrt{\frac{N-t}{N}} |B\rangle \quad (13)$$

Proof. By Definition 4, the uniform superposition state $|U\rangle$ is given by:

$$|U\rangle = \frac{1}{\sqrt{N}} \sum_{x \in U} |x\rangle \quad (14)$$

Splitting the sum into tagged ($x \in A$) and untagged ($x \in B$) elements:

$$|U\rangle = \frac{1}{\sqrt{N}} \left(\sum_{x \in A} |x\rangle + \sum_{x \in B} |x\rangle \right) \quad (15)$$

From Definition 4, we have $\sum_{x \in A} |x\rangle = \sqrt{t} |A\rangle$ and $\sum_{x \in B} |x\rangle = \sqrt{N-t} |B\rangle$. Substituting these:

$$|U\rangle = \frac{1}{\sqrt{N}} \left(\sqrt{t} |A\rangle + \sqrt{N-t} |B\rangle \right) = \sqrt{\frac{t}{N}} |A\rangle + \sqrt{\frac{N-t}{N}} |B\rangle \quad (16)$$

□

3 Grover's Algorithm

Definition 11 (Grover Iterate). The Grover iterate G (also called the Grover operator) is defined as:

$$G = H^{\otimes n} (2|0\rangle\langle 0| - I) H^{\otimes n} Z_f \quad (17)$$

where Z_f is the function oracle and $H^{\otimes n} (2|0\rangle\langle 0| - I) H^{\otimes n}$ is the reflection about the uniform superposition state $|U\rangle$ (as shown in Corollary 9).

Grover's algorithm consists of starting with the uniform superposition state $|U\rangle$ and repeatedly applying the Grover iterate G . Geometrically, each application of G rotates the state vector in the 2D plane spanned by $|A\rangle$ and $|B\rangle$ towards the target state $|A\rangle$.

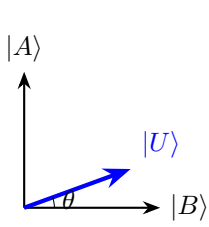


Figure 1: Uniform state $|U\rangle$

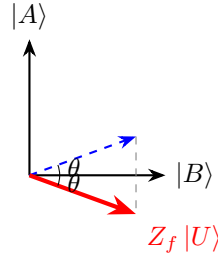


Figure 2: Action of Z_f

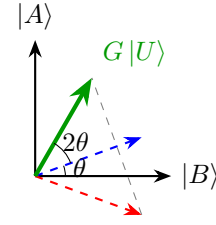


Figure 3: Action of G

Definition 12 (Grover Angle). Define the angle θ such that:

$$\sin \theta = \sqrt{\frac{t}{N}} \quad \text{and} \quad \cos \theta = \sqrt{\frac{N-t}{N}} \quad (18)$$

This allows us to write the uniform superposition state (from Proposition 10) as:

$$|U\rangle = \sin \theta |A\rangle + \cos \theta |B\rangle \quad (19)$$

where θ is the angle between the initial state $|U\rangle$ and the *bad* state superposition $|B\rangle$.

Proposition 13 (G as a Rotation). The Grover iterate $G = (2|U\rangle\langle U| - I)Z_f$ is a rotation by 2θ in the plane spanned by $\{|A\rangle, |B\rangle\}$.

Proof. The operator G is the composition of two reflections:

1. Z_f is a reflection about the axis $|B\rangle$. For any state $|\psi\rangle = \alpha|A\rangle + \beta|B\rangle$, $Z_f|\psi\rangle = -\alpha|A\rangle + \beta|B\rangle$, which negates the component orthogonal to $|B\rangle$ ($x \in A \implies f(x) = 1 \implies Z_f|x\rangle = -|x\rangle$).
2. $2|U\rangle\langle U| - I$ is a reflection about the axis $|U\rangle$ (as shown in Lemma 8).

From geometry, the composition of two reflections about lines intersecting at an angle θ is a rotation by 2θ in the direction from the first reflection axis to the second. Since the angle between $|B\rangle$ and $|U\rangle$ is θ , the result is a rotation by 2θ towards $|A\rangle$. \square

Definition 14 (State After k Iterations). Let $|\varphi_k\rangle$ be the state obtained after k applications of the Grover iterate G starting from the initial uniform superposition $|U\rangle$:

$$|\varphi_k\rangle = G^k |U\rangle \quad (20)$$

From Proposition 13, since G is a rotation by 2θ in the plane $\{|A\rangle, |B\rangle\}$ and the starting state $|U\rangle$ is at an angle θ with the axis $|B\rangle$, the state after k iterations is:

$$|\varphi_k\rangle = \sin((2k+1)\theta) |A\rangle + \cos((2k+1)\theta) |B\rangle \quad (21)$$

Definition 15 (Optimal Iterations). To maximize the probability of measuring a *good* state $|A\rangle$, we want $(2k+1)\theta \approx \pi/2$. The optimal number of Grover iterations k^* is given by the closest integer:

$$k^* = \left\lfloor \frac{\pi}{4\theta} - \frac{1}{2} \right\rfloor \quad (22)$$

For small $\theta \approx \sin \theta = \sqrt{t/N}$, this yields $k^* \approx \frac{\pi}{4} \sqrt{\frac{N}{t}}$.

Proposition 16 (Success Probability). *Using k^* Grover iterations ensures the probability of measuring a good state $x \in A$ is at least $1/2$.*

Proof. If $\sin^2 \theta \geq 1/2$, the initial uniform superposition state $|U\rangle$ already has a high probability of yielding a tagged element. In this case, $k^* = 0$, and we succeed with probability $\sin^2 \theta \geq 1/2$.

Otherwise, we have $\sin^2 \theta < 1/2$. The optimal number of iterations k^* is chosen such that the final angle $(2k^* + 1)\theta$ is as close as possible to $\pi/2$. Since k^* is the closest integer to $\frac{\pi}{4\theta} - \frac{1}{2}$, the distance from the target angle is at most θ :

$$\left| k^* - \left(\frac{\pi}{4\theta} - \frac{1}{2} \right) \right| \leq \frac{1}{2} \implies |(2k^* + 1)\theta - \pi/2| \leq \theta \quad (23)$$

The probability of failure (measuring a state in B) is given by:

$$P(B) = \cos^2((2k^* + 1)\theta) = \sin^2\left(\frac{\pi}{2} - (2k^* + 1)\theta\right) \leq \sin^2 \theta \quad (24)$$

Since we assumed $\sin^2 \theta < 1/2$, the failure probability is bounded by $1/2$, which implies the success probability is $P(A) = 1 - P(B) \geq 1 - \sin^2 \theta > 1/2$. \square

4 Amplitude Amplification

Definition 17 (Amplitude Amplification Problem). Given a unitary operator A such that:

$$A|0^n\rangle = \sqrt{p}| \psi_0\rangle + \sqrt{1-p}| \psi_1\rangle \quad (25)$$

where:

- $|\psi_0\rangle$ is the superposition of *good* states: $|\psi_0\rangle = \sum_i \alpha_i |x_i\rangle$, where $\alpha_i \neq 0$ if and only if $f(x_i) = 1$.
- $|\psi_1\rangle$ is the superposition of *bad* states: $|\psi_1\rangle = \sum_i \beta_i |x_i\rangle$, where $\beta_i \neq 0$ if and only if $f(x_i) = 0$.

The task is to amplify the probability of $|\psi_0\rangle$ (initially p) while keeping the relative amplitudes α_i and β_i within each superposition intact.

Remark 18 (Grover Operator in Amplitude Amplification). Given the unitary operator A and its adjoint A^\dagger , we can construct the generalized Grover iterate G :

$$G = A(2|0\rangle\langle 0| - I)A^\dagger Z_f \quad (26)$$

where Z_f is the reflection about the *good* state subspace (as defined before). The geometric analysis remains the same, with $\sin^2 \theta = p$, allowing us to use Grover's algorithm to amplify the success probability from p to at least $1/2$ (or higher).

4.1 Exact Amplitude Amplification

Proposition 19 (Reaching Success Probability 1). *By introducing an auxiliary qubit and a rotation operator U , it is possible to reach the good state $|\psi_0\rangle$ with probability exactly 1.*

Proof. Introduce an auxiliary qubit and define the rotation operator U :

$$U = \begin{bmatrix} \cos \varphi & -\sin \varphi \\ \sin \varphi & \cos \varphi \end{bmatrix} \quad (27)$$

We define a new preparation operator $A^* = A \otimes U$ acting on $n + 1$ qubits. Its action on the state $|0^{n+1}\rangle$ is:

$$A^* |0^{n+1}\rangle = \left(\sqrt{p} |\psi_0\rangle + \sqrt{1-p} |\psi_1\rangle \right) \otimes (\cos \varphi |0\rangle + \sin \varphi |1\rangle) \quad (28)$$

Expanding this yields:

$$A^* |0^{n+1}\rangle = \underbrace{\sqrt{p} \cos \varphi |\psi_0\rangle |0\rangle}_{\text{Good State}} + \underbrace{\left(\sqrt{p} \sin \varphi |\psi_0\rangle |1\rangle + \sqrt{1-p} \cos \varphi |\psi_1\rangle |0\rangle + \sqrt{1-p} \sin \varphi |\psi_1\rangle |1\rangle \right)}_{\text{Bad States}} \quad (29)$$

Define a new oracle $f^*(x, y)$ (where $x \in \{0, 1\}^n$ and $y \in \{0, 1\}$) such that:

$$f^*(x, y) = \begin{cases} f(x) & \text{if } y = 0 \\ 0 & \text{if } y = 1 \end{cases} \quad (30)$$

Under this oracle, the only tagged states are of the form $|x\rangle |0\rangle$ where $f(x) = 1$. The initial success amplitude is $\sin \theta^* = \sqrt{p} \cos \varphi$.

To ensure we reach probability 1, we first choose the number of iterations k^* as:

$$k^* = \left\lceil \frac{\pi}{4\theta} - \frac{1}{2} \right\rceil \quad (31)$$

where $\sin \theta = \sqrt{p}$. This choice of k^* ensures that:

$$2k^* + 1 \geq \frac{\pi}{2\theta} \implies \theta^* = \frac{\pi}{2(2k^* + 1)} \leq \theta \quad (32)$$

Since $\theta^* \leq \theta$, we have $\sin \theta^* \leq \sin \theta = \sqrt{p}$. Thus, we can always find a rotation angle φ such that:

$$\sqrt{p} \cos \varphi = \sin \theta^* \implies \cos \varphi = \frac{\sin \theta^*}{\sqrt{p}} \leq 1 \quad (33)$$

With this choice of φ , the state after k^* iterations of the Grover operator G^* (constructed using A^* and f^*) will be:

$$|\varphi_{k^*}\rangle = \sin((2k^* + 1)\theta^*) |\text{Good}\rangle + \cos((2k^* + 1)\theta^*) |\text{Bad}\rangle \quad (34)$$

Substituting $(2k^* + 1)\theta^* = \pi/2$ yields:

$$|\varphi_{k^*}\rangle = \sin(\pi/2) |\text{Good}\rangle + \cos(\pi/2) |\text{Bad}\rangle = |\text{Good}\rangle \quad (35)$$

Thus, we reach the target state $|\psi_0\rangle |0\rangle$ with probability exactly 1. \square